



UNIVERSITI PUTRA MALAYSIA

**DATA LINK CONTROL LAYER PERFORMANCE FOR WIRELESS
ATM NETWORKS**

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**DATA LINK CONTROL LAYER PERFORMANCE FOR WIRELESS ATM
NETWORKS**

By

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**Thesis Submitted in Fulfilment of the Requirements for the Degree of Master of
Science in the Faculty of Engineering
Universiti Putra Malaysia**

April 2000



Dedicated to

My Parents, Salma and Khalifa Yaklaf, and to my wife Aziza and my children (Aimen, Monia and Tasnim) for their patience and encouragement throughout this work.

Abstract of thesis presented to the Senate of Universiti Putra Malaysia in fulfilment of the requirements for the degree of Master of Science.

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April 2000

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The growing demand for ATM-based technology and recent proliferation of wireless access technologies have motivated researchers to examine the feasibility of extending the ATM paradigm from the wireline to the wireless domain and create a new research area known as Wireless ATM (WATM) (Toh, 1997). Dealing with lossy wireless links, characterized by limited bandwidth and high, bursty error rates, breaks the main assumption of conventional ATM systems, which is that of using no errors per links.

Therefore, WATM systems must provide a transparent mechanism to ensure reliable end-to-end data transmission over the wireless portion of the network. The identification of a wireless-specific data link control layer (W-DLC), sitting between the traditional ATM layer and a wireless-specific medium access control layer (W-MAC), is the responsible entity for guaranteeing the quality of service (QOS) requested by individual ATM-based virtual connections. Thus the main focus is to investigate the performance of DLC protocol for ABR traffic over wireless ATM network. Retransmissions are only required for non-real time traffic and are

implemented using a Go-Back-N and Selective Repeat (SR) ARQ (Lin et al., 1984), (Schwartz, 1987).

Wireless channels are usually time-varying and the channel bit error rates vary as the surrounding environment changes. Since these factors put in jeopardy the performance of the DLC protocol and higher layer end-to-end protocol at large, additional link-level mechanisms are added to provide reliability over impaired radio links. The DLC protocol implementation represents an attempt to achieve these goals under the strict constraints imposed by impaired wireless links.

This thesis studies and compares the two Automatic Repeat Request (ARQ) protocols, i.e., Go-Back-N (GBN) and Selective Repeat (SR) ARQ and analyse them for variable packet size and fixed packet size (WATM packet) by using C programming for simulation.

The results show that the performance of SR ARQ is better than the GBN ARQ for variable packet size. The results also show that SR ARQ protocol has better performance than GBN ARQ in terms of error detection for fixed WATM packet in the range of 50 – 70 bytes, which is the WATM packet range (ATM Forum, 1997).

Abstrak tesis yang dikemukakan kepada Senat Universiti Putra Malaysia
sebagai memenuhi keperluan untuk ijazah Master Sains.

**PRESTASI LAPISAN KAWALAN PAUTAN DATA UNTUK RANGKAIAN
ATM WAYARLES**

Oleh

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Permintaan yang semakin meningkat untuk teknologi berasaskan ATM dan pertumbuhan pesat teknologi capaian wayarles sekarang telah menggalakkan penyelidik untuk mengkaji kebolehlaksanaan mengembangkan paradigma ATM daripada jalur wayar kepada domain wayarles dan mencipta satu ruang penyelidikan baru yang dikenali sebagai ATM Wayarles (WATM) (Toh, 1997). Berurusan dengan hubungan wayarles “terhilang”, disebabkan dibentuk oleh lebar jalur yang terhad dan kadar ralat yang tinggi dan bertaburan, telah memecahkan jangkaan utama sistem ATM konvensional yang melibatkan penggunaan setiap hubungan tanpa ralat.

Oleh yang demikian, sistem WATM mesti menyediakan satu mekanisme lutsinar untuk memastikan penghantaran data hujung-ke-hujung yang diyakini menerusi bahagian wayarles sesuatu rangkaian. Pengenalpastian satu lapisan kawalan data berkhususkan-wayarles (W-DLC), yang berada di antara lapisan ATM tradisional dan lapisan kawalan capaian medium berkhususkan wayarles (W-MAC), adalah entiti yang bertanggungjawab untuk menjamin kualiti perkhidmatan (QoS) yang dikehendaki oleh sambungan maya individu berasaskan-ATM. Oleh itu, fokus

kami adalah menyelidik prestasi protokol DLC kami untuk trafik ABR menerusi rangkaian ATM wayarles. Penghantaran semula hanya dikehendaki untuk trafik bukan masa nyata dan dilaksanakan menggunakan satu protocol ARQ Balik-ke-N dan Ulangan Terpilih (SR) (Lin et al., 1984), (Schwartz, 1987).

Saluran-saluran wayarles biasanya berubah mengikut masa dan kadar ralat bit saluran berubah mengikut perubahan persekitaran. Oleh kerana faktor-faktor ini sangat mengancam prestasi protokol DLC dan protokol hujung-ke-hujung lapisan tertinggi, mekanisma-mekanisma peringkat-pautan tambahan dimasukkan untuk memberi kebolehpercayaan ke atas pautan radio yang terjejas. Pelaksanaan protokol DLC kami adalah satu percubaan untuk mencapai matlamat-matlamat tersebut dalam kekangan ketat yang dikenakan oleh hubungan wayarles terjejas.

Dalam tesis ini, kami menyelidik dan membanding dua Protokol Permohonan Ulangan Automatik (ARQ) yakni Balik-ke-N (GBN) dan Ulangan Terpilih (SR) serta menganalisis untuk saiz paket bolehubah dan saiz paket tetap (paket WATM) dengan menggunakan pengaturcaraan C di dalam simulasi kami.

Hasilnya menunjukkan prestasi ARQ SR adalah lebih baik daripada ARQ GBN untuk saiz paket bolehubah. Keputusan itu juga menunjukkan protokol ARQ SR mempunyai prestasi yang lebih baik dari ARQ GBN dari segi pengesanan ralat untuk WATM tetap dalam julat 50-70 bait, seperti mana yang dicadangkan oleh (ATM Forum, 1997).

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I certify that an Examination Committee met on 28th April 2000 to conduct the final examination of Sliman Ka. A. Yaklaf on his Master of Science thesis entitled “ Data Link Control Layer Performance for Wireless ATM Networks” in accordance with Universiti Pertanian Malaysia (Higher Degree) Act 1980 and Universiti Pertanian Malaysia (Higher Degree) Regulations 1981. The Committee recommends that the candidate be awarded the relevant degree. Members of the Examination Committee are as follows:

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DECLARATION

I hereby declare that the thesis is based on my original work except for quotations and citations which have been duly acknowledged. I also declare that it has not been previously or concurrently submitted for any other degree at UPM or other institutions.



SLIMAN KA. A. YAKLAF

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TABLE OF CONTENTS

	Page
DEDICATION.....	2
ABSTRACT.....	3
ABSTRAK.....	5
ACKNOWLEDGEMENT.....	7
APPROVAL SHEETS.....	8
DECLARATION FORM.....	10
LIST OF TABLES.....	15
LIST OF FIGURES.....	16
LIST OF ABBREVIATIONS	19
CHAPTER	
I INTRODUCTION.....	
Background.....	21
Motivation and History.....	22
Wireless ATM Protocol Architecture.....	23
Data Link Layer Protocols for Wireless ATM.....	23
Research Objectives.....	25
Contribution of The Study.....	25
Thesis Organisation	26
II LITERATURE REVIEW	
Wireless ATM (rational and challenges)	28
Wireless ATM Specifications.....	32
Radio Access Protocols	32
Radio Physical Layer	33
Medium Access Control	34
Data Link Control	35
Data Link Layer Performance	36
Fame Handling	37
Flow Control	37
Error Control	38
Error Detection	42
Link Management	43
Brief Overview of Wireless ATM Systems.....	43
Seamless Wireless ATM Networks	
(SWAN) - Bell Laboratories	43
Broadband Ad-Hoc Wireless ATM LAN	
(BAHAMA) – Bell Laboratories	44
Wireless ATM Networks (WATMnet) –	
NEC C&C Research Laboratories	45
Cambridge – Olivetti Research Laboratories	45
Carleton University	46



	Advanced Communication Technologies and Services (ACTS) Magic Wireless ATM Networks Demonstrator (ACTS Magic WAND)	47
	Asymmetric Reliable Mobile Access In Link-Layer (AIRMAIL) – AT & T Bell Lab.....	47
	Need for Data Link Control Protocol	48
	Brief Overview of Data Link Control in Previous Wireless ATM Systems	49
	SWAN	50
	BAHAMA	51
	WATMnet	51
	Cambridge – Olivetti Laboratories	52
	Carleton University	53
	ACTS Magic WAND	53
	AIRMAIL	54
	Data Link Control in Packet Radio Networks	55
	The implementation of Data Link Control for Wireless ATM Networks	56
	Conclusions	57
III	METHODOLOGY	59
	Simulation Model of GBN and SR ARQ for Wireless ATM Networks.....	59
	Theoretical Model	59
	Wireless ATM Packet Format	60
	The Effect of Channel Noise	62
	Window Size	62
	Adaptive Packet Length	62
	Objective of the Simulation	63
	Description of the Simulation Model	64
	Data Source Parameters	64
	System Parameters.....	65
	Designing the Simulator	65
	Operation of the ARQ Protocols	66
	Simulation Model of GBN ARQ Protocol	67
	Input Parameters	67
	Output Parameters	67
	Transmitter Algorithm	68
	Receiver Algorithm	68
	Simulation Model of SR ARQ Protocol	70
	Transmitter Algorithm	71
	Receiver Algorithm	71
	Summary	73
IV	RESULTS AND DISCUSSION	75
	Introduction	75
	Evaluation and Performance Comparison of Go-Back-N Protocol for Different Packet Sizes	76
	Packet Sizes (55 – 199 bytes)	76
	Packet Sizes (200 – 500 bytes)	80

	Fixed Packet Size (54 bytes)	81
	Performance Comparison of Go-Back-N for Different Packet Sizes	85
	Effect of Different Error Rates on Good Packet Received for (55 – 199 bytes)	86
	Effect of Different Error Rates on Throughput for (200 – 500 bytes)	87
	Effect of Different Error Rates on Good Packet Received for Fixed Packet Size (54 bytes)	88
	Evaluation and Performance Comparison of Selective Repeat ARQ Protocol for Different Packet Sizes	89
	Packet Sizes (55 – 199 bytes)	89
	Packet Sizes (200 – 500 bytes)	90
	Fixed Packet Size (54 bytes)	91
	Performance Comparison of Selective Repeat ARQ for Different Packet Sizes	92
	Performance Comparison of Go-Back-N and Selective Repeat ARQ Protocols for Different packet Sizes	93
	Packet Sizes (55 – 199 bytes)	93
	Packet Sizes (200 – 500 bytes)	93
	Fixed Packet Size (54 bytes)	95
	Different Packet Sizes	95
	Effect of (15 – 25%) Error Rates on Go-Back-N and Selective Repeat ARQ Performance	96
	Summary	99
V	CONCLUSION AND FUTURE WORKS	101
	Conclusions	101
	Future works	103
	REFERENCES	104
	APPENDICES	107
A-1	The effect of (15 – 35%) packet losses and check-sum error on throughput of GBN Protocol for (55 – 199 bytes) packet sizes (Figures)	108
A-2	The effect of (10 – 15%) packet losses and check-sum error on throughput of GBN protocol for (200 – 500 bytes) packet sizes (Figures)	111
A-3	The effect of (15 – 35%) packet losses and check-sum error on throughput of GBN protocol for fixed packet (54 bytes) (Figures)	112
A-4	(Tables)	114
	Table 3: Performance Comparison for GBN and SR ARQ for Packet Sizes (55 – 199 bytes)	114

Table 4: Performance Comparison for GBN and SR ARQ for Packet Sizes (200 - 500 bytes)	114
Table 5: Performance Comparison for GBN and SR ARQ for Fixed Packet (54) bytes	115
BIODATA OF AUTHOR	116

LIST OF TABLES

Table		Page
1	Simulation Parameters of GBN ARQ for Variable Packet Sizes	76
2	Simulation parameters of GBN ARQ for Fixed Packet Size	82
3	Performance Comparison for GBN and SR ARQ for Packet Sizes (55 – 199 bytes)	114
4	Performance Comparison for GBN and SR ARQ for Packet Sizes (200 – 500 bytes)	114
5	Performance Comparison for GBN and SR ARQ for Fixed Packet Size (54 bytes)	115

LIST OF FIGURES

Figure		Page
1	Typical Wireless Enhanced ATM Protocol Stack	31
2	Wireless ATM Protocol Stack	33
3	TDMA/TDD Frame Format for Use in Wireless ATM	35
4	Logical Representation of the Data Link Control Layer	36
5	Operation of Go-Back-N ARQ Protocol	40
6	Operation of Selective Repeat ARQ Protocol	42
7	General Network Model for Simulation	60
8	Standard ATM Cell Format	60
9	WATM Cell Format	61
10	Transmitter Flow Chart for GBN ARQ Protocol Process	69
11	Receiver Flow Chart for GBN ARQ Protocol Process	70
12	Transmitter flow Chart for SR ARQ Protocol Process	72
13	Receiver Flow Chart for SR ARQ Protocol Process	73
14	The effect of timeout on GBN throughput with 1% packet losses and check-sum error for (55 – 199 bytes)	77
15	The effect of timeout on GBN throughput with 5% packet losses and check-sum error for (55 – 199 bytes)	78
16	The effect of timeout on GBN throughput with 10% packet losses and check-sum error for (55 – 199 bytes)	79
17	The effect of timeout on GBN throughput with 1% packet losses and check-sum error for (200 – 500 bytes)	80
18	The effect of timeout on GBN throughput with 5% packet losses and check-sum error for (200 – 500 bytes)	81
19	The effect of timeout on GBN throughput with 1% packet losses and check-sum error for (54 bytes)	83

20	The effect of timeout on GBN throughput with 5% packet losses and check-sum error for (54 bytes)	83
21	The effect of timeout on GBN throughput with 10% packet losses and check-sum error for (54 bytes)	84
22	The effect of different packet losses and check-sum error rates on good packets received of GBN protocol for different packet sizes	86
23	Effect of different error rates on good packet received of GBN protocol for (55 – 199 bytes)	87
24	Effect of different error rates on throughput of GBN for (200 - 500 bytes)	87
25	Effect of different Error Rates on good packet received of GBN protocol for (54 bytes)	88
26	Effect of different packet losses and check-sum error rates on SR throughput for (55 – 199 bytes)	89
27	Effect of different packet losses and check-sum error rates on SR throughput for (200 - 500 bytes)	90
28	Effect of different packet losses and check-sum error rates on SR throughput for (54 bytes)	91
29	Effect of different packet losses and check-sum error rates on SR throughput for different packet sizes	92
30	Effect of different packet losses and check-sum error rates on GBN and SR throughput performance for (55 – 199 bytes).....	94
31	Effect of different packet losses and check-sum error rate on GBN and SR throughput performance for (200 – 500 bytes)	94
32	Effect of different packet losses and check-sum error rates on GBN and SR throughput performance for (54 bytes)	95
33	Effect of different packet losses and check-sum error rates on GBN and SR throughput performance for different packet sizes	96
34	Effect of packet size on good packet received with 15% packet losses and check-sum error for GBN and SR ARQ	97
35	Effect of packet size on good packets received with 20% packet losses and check-sum error for GBN and SR ARQ	98

36	Effect of packet size on good packets received with 25% packet losses and check-sum error for GBN and SR ARQ	98
37	The effect of timeout on GBN throughput with 15% packet losses and check-sum error for (55 – 199 bytes)	108
38	The effect of timeout on GBN throughput with 20% packet losses and check-sum error for (155 – 199 bytes).....	108
39	The effect of timeout on GBN throughput with 25% packet losses and check-sum error for (55 – 199 bytes).....	109
40	The effect of timeout on GBN throughput with 30% packet losses and check-sum error for (55 – 199 bytes).....	109
41	The effect of timeout on GBN throughput with 35% packet losses and check-sum error for (55 – 199 bytes).....	110
42	The effect of timeout on GBN throughput with 10% packet losses and check-sum error for (200 – 500 bytes).....	111
43	The effect of timeout on GBN throughput with 15% packet losses and check-sum error for (200 – 500 bytes).....	111
44	The effect of timeout on GBN throughput with 15% packet losses and check-sum error for (54 bytes)	112
45	The effect of timeout on GBN throughput with 20% packet losses and check-sum error for (54 bytes)	112
46	The effect of timeout on GBN throughput with 25% packet losses and check-sum error for (54 bytes)	113
47	The effect of timeout on GBN throughput with 30% packet losses and check-sum error for (54 bytes)	113

LIST OF ABBREVIATION

AAL	ATM Adaptation Layer
ABR	Available Bit Rate
ACK	Acknowledgement
ACTS	Advanced Communication Technologies and Services
AIRMAIL	Asymmetric Reliable Mobile Access In Link-Layer
AP	Access Point
ARQ	Automatic Repeat Request
ATM	Asynchronous Transfer Mode
ATMLC	ATM Link Control
BAHAMA	Broadband Ad-Hoc Wireless ATM LAN
BER	Bit Error Rate
B-R	Base-to-Remote
CBR	Constant Bit Rate
CITR	Canadian Institute of Telecommunication Research
CLP	Cell Loss Priority
CRC	Cyclic Redundancy Check
CSN	Cell Sequence Number
DLC	Data Link Control
DQRUMA	Distributed Queuing Request Up-date Multiple Access
ETSI	European Telecommunications Standards Institute
FEC	Forward Error Correction
FIFO	First-In-First-Out
GBN	Go-Back-N
HI	Hand-off-Indicator
IP	Internet Protocol
LAN	Local Area Network
MAC	Medium Access Control
MAX-SEQ	Maximum-Sequence
MT	Mobile Terminal
NAK	Negative Acknowledgement
PBS	Portable Base Station
PC	Personal Computer
PDU	Protocol Data Unit
PHY	Physical Layer
PTI	Payload Type Indicator
QoS	Quality of Service
R-B	Remote-to-Base
RRC	Radio Resource Control
SDLC	Synchronous Data Link Control
SN	Sequence Number
SR	Selective Repeat
SREJ	Selective Reject
SW	Stop-and-Wait
SWAN	Seamless Wireless ATM Network
TCP	Transmission Control Protocol
TDD	Time Division Duplex

TDMA	Time Division Multiplex Access
ToE	Time of Expiry
UBR	Unspecified Bit Rate
UTP	Unshielded Twisted Pair
VBR	Variable Bit Rate
VC	Virtual Channel
VCI	Virtual Channel Identifier
VCs	Virtual Connections
VP	Virtual Path
VPI	Virtual Path Identifier
WAN	Wide Area Network
WAND	Wireless ATM Network Demonstrator
WATM	Wireless ATM
WATMnet	Wireless ATM network
W-DLC	Wireless Specific Data Link Control
W-MAC	Wireless Specific Medium Access Control

CHAPTER I

INTRODUCTION

Background

The last few years have seen an explosion in wireline broadband networking technologies. This has been driven by an increased user demand for video teleconferencing, image applications, world-wide web access and other multimedia applications as well as advancements in key enabling technologies such as high-speed digital transmission using optical fibre, digital signal processing and high-speed integrated circuits. In this time Asynchronous Transfer Mode (ATM) has emerged as the front runner in integrated telecommunication technologies. ATM provides high-speed transfer, integration of traffic types, flexible bandwidth allocation and service type selection for a range of applications, efficient multiplexing of data from bursty data/multimedia sources and simplified network management. It is rapidly becoming a world-wide standard and has moved from concept to reality in the space of a few years.

Wireless personal communication networks have also emerged as an important field of activity in telecommunications. This surge of interest is due to several factors such as the increased availability of wireless personal computing, entertainment and communication devices, and liberalisation of spectrum allocation procedures and advances in digital signal processing and radio modem technologies. While these systems have initially focused only on voice and primitive packet data

applications, it is recognised that they will have to evolve toward supporting a wider range of applications involving video and multimedia. The increased dependence on networking for business, recreation and communications, the growing demand for multimedia applications together with a human desire for mobility and freedom from office-only or home-only computing constraints makes a strong argument for wireless integrated networks.

Motivation and History

The growing interest in wireless integrated networks has motivated several researchers to examine the feasibility of extending the ATM paradigm from the wireline to the wireless domain, and create new research area known as wireless ATM (WATM) (Toh, 1997). ATM was designed for a time-invariant, reliable medium where bandwidth is not a significant constraint. However, the wireless channel is usually time varying with a high bit error rate and limited bandwidth.

The ATM protocol provides no built-in mechanism to recover from errors or cell losses due to problems with the underlying transport medium. Thus one of the requirements for making wireless ATM a reality is a method of compensating for the low reliability of the wireless channel. The obvious solution to this problem is to create a protocol at the data link level that handles the error recovery and control over the wireless channel and provides the ATM layer with a more reliable transport medium.

Wireless ATM Protocol Architecture

Wireless ATM adopts ATM to provide the data communications services so the overall architecture is based on the ATM protocol stack. A wireless ATM protocol architecture is currently proposed by ATM Forum. The WATM items are divided into two distinct parts:

- Mobile ATM
- Radio Access Layer

Mobile ATM deals with the higher layer control and signalling functions needed to support mobility. These control and signalling include handover, location management, routing, addressing, and traffic management.

Radio Access layer is responsible for the radio link protocols for wireless ATM access. Radio Access Layer consists of Physical Layer (PHY), Medium Access Control Layer (MAC), Data Link Layer (DLC), and Radio Resource Control (RRC).

Data Link Control Layer Protocols for Wireless ATM

The DLC layer is responsible for providing services to the ATM layer. Mitigating the effect of radio channel errors should be done in this layer before cells are sent to the ATM layer. Since end-to-end ATM performance is sensitive to cell loss, powerful error control procedures are required for the WATM radio access segment. In order to fulfil this requirement, Automatic Repeat Request (ARQ) schemes are used to improve the performance of communication. So, error

detection/retransmission protocols and forward error correction methods are recommended.

The three main protocols used in ARQ schemes are stop-and-wait (SW), go-back-N (GBN), and selective repeat (SR). The SW protocol is inefficient due to the idle time spent waiting for the receiver to acknowledge each transmission. GBN and SR schemes make use of sliding window principle in which the window slides along as sequence numbers of packets in a transmit window are acknowledged, enabling more packets to be sent. In GBN scheme, if an error is detected in a packet, the receiver discards all subsequent packets, until an uncorrupted version of this packet is received. In GBN protocol, the loss of a packet is detected only after the next packet has been correctly received. This requires a continuous stream of packets being ready to send; otherwise unacceptably long delays are experienced while the receiver waits for the next packet which is transmitted by the sender, and is stopped when an acknowledgement (ACK) indicating its correct reception is received. If an ACK is not received before its time-out interval expires, the packet is retransmitted.

In selective repeat (SR) protocol, if a corrupted packet is received, it is discarded, but this does not have any effect on the acceptance of subsequent packets. This protocol uses the channel bandwidth efficiently, but requires more buffering at the receiver than Go-Back-N ARQ technique. The SR scheme uses two ways to retransmit:

- Implicit retransmission where the transmitter determines from the sequence of ACK frames that a frame has been lost.